

# MA633L-Numerical Analysis

Lecture 26 : Numerical Linear Algebra - Norms

**Panchatcharam Mariappan<sup>1</sup>**

<sup>1</sup>Associate Professor  
Department of Mathematics and Statistics  
IIT Tirupati, Tirupati

**March 10, 2025**





# Numerical Weather Prediction

# Numerical Weather Prediction

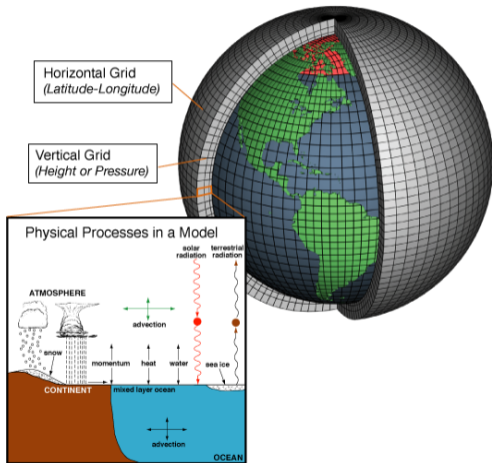


Figure 1: Atmospheric Model Grids, Source: Wikipedia

# Numerical Weather Prediction

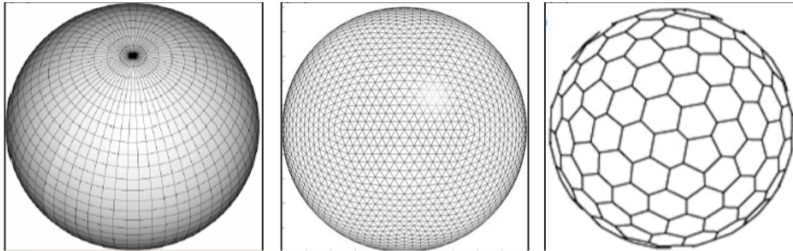
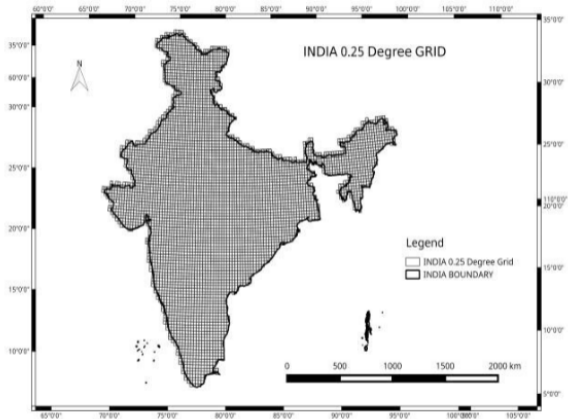


Figure 2: Atmospheric Model Grids Other Pattern, Source: Intech Open

# Numerical Weather Prediction



**Figure 3:** Atmospheric Model Grids: India, Source: Forest Fire Locations in India their spatio temporal patterns

# Numerical Weather Prediction



- Weather Prediction of India for the next 10 days
- Area of India is 3.2 million sq.km
- To predict the weather pattern, you should model the atmosphere from sea level to 20 km
- Assume that, we make prediction of weather at each cubical grid with each cube measuring 0.1 km on each side.
- Compute  $F(\text{lattitude, longitude, elevation, time}) = \text{Temperature or Pressure or Rain or ...}$

# Numerical Weather Prediction



- Total number of points =  $3.2 \times 10^6 \times 20 \times 10^3 = 6.4 \times 10^{10}$
- If these nodes are connected to each other and influences, the size of the matrix is  $10^{10}$

# Gaussian Elimination



- To solve the NWP matrix, of size  $10^{10}$  using GFLOP machine, it requires,  $10^{30}/10^9 = 10^{21}$  seconds, that is,  $10^{13}$  years to compute, but age of earth is  $4.5 \times 10^9$  years.
- For a tera flops machine (1 TFLOPS =  $10^{12}$  FLOPS), it requires,  $10^{10}$  years.
- For a peta flops machine (1 PFLOPS =  $10^{15}$  FLOPS), it requires,  $3.16 \times 10^7$  years.
- Researchers use HPC, Supercomputers and sparse linear solvers





# Sparse Matrix and Iterative Schemes

# Sparse Matrix



Note: there is no strict or standard definition regarding the sparse matrix, but the following is commonly used

## Definition 1 (Sparse Matrix)

A sparse matrix is a matrix in which most of the elements are zero. A common criterion is that the number of zeros is roughly equal to the number of rows or column.

## Definition 2 (Sparsity)

The sparsity of matrix can be represented as

$$\text{Sparsity} = \frac{\text{Number of Zero Elements}}{\text{Total Number of Elements}}$$

Typically, a matrix is considered sparse if sparsity is  $> 0.5$

# Iterative Schemes

The underlying principles behind iterative methods to solve  $Ax = b$  are as follows:

1. Guess any  $x^{(0)}$
2. Compute the residual  $r^{(0)} = b - Ax^{(0)}$
3. Compute the residual norm, that is  $\|r^{(0)}\|$
4. Use an algorithm to compute  $x^{(1)}$  involving  $A, b$  and  $x^{(0)}$
5. Recompute the residual norm  $\|r^{(1)}\| = \|b - Ax^{(1)}\|$
6. For given  $x^{(i)}$ , iterate this process until you get the residual as zero or an equivalent condition



# Vector Norms

# Norms



## Definition 3 (Vector Norms)

A vector norm, usually denoted by  $\|\cdot\|$ , is a function from a vector space  $V$  to the set of nonnegative real numbers that obeys the following three postulates.

$\|\cdot\| : V \rightarrow \mathbb{R}_+$  such that

$$\|\mathbf{x}\| > 0 \quad \text{if } \mathbf{x} \neq 0, \mathbf{x} \in V$$

$$\|\alpha\mathbf{x}\| = |\alpha|\|\mathbf{x}\| \quad \text{if } \alpha \in \mathbb{R}, \mathbf{x} \in V$$

$$\|\mathbf{x} + \mathbf{y}\| \leq \|\mathbf{x}\| + \|\mathbf{y}\| \quad \text{if } \mathbf{x}, \mathbf{y} \in V$$

# Vector Norms



The most familiar norm on  $\mathbb{R}^n$  is the Euclidean norm defined by

$$\|\mathbf{x}\|_2 = \left( \sum_{i=1}^n |x_i|^2 \right)^{1/2}, \quad \text{where } \mathbf{x} = (x_1, x_2, \dots, x_n) \quad (1)$$

This is the norm that corresponds to usual concept of length. The other norms are also used. The simplest and easiest norm is  $\ell_\infty$ -norm

$$\|\mathbf{x}\|_\infty = \max_{0 \leq i \leq n} |x_i| \quad (2)$$

# Vector Norms



The third important norm is  $\ell_1$ -norm

$$\|\mathbf{x}\|_1 = \sum_{i=1}^n |x_i| \quad (3)$$

## Example 4

Compute the  $\|\cdot\|_1$ ,  $\|\cdot\|_2$  and  $\|\cdot\|_\infty$  norms of the following vectors.

$x = (4, 4, -4, 4)$   $v = (0, 2, 2, 2)$ ,  $w = (8, 0, 0, 0)$

	$\ \cdot\ _1$	$\ \cdot\ _2$	$\ \cdot\ _\infty$
$x$	16	8	4
$v$	6	$\sqrt{12}$	2
$w$	8	8	8

# Vector Norms



The  $p$ - norm or  $\ell_p$  norm of a vector is defined by

$$\|\mathbf{x}\|_p = \left( \sum_{i=1}^n |x_i|^p \right)^{1/p}, \quad \text{where } \mathbf{x} = (x_1, x_2, \dots, x_n) \quad (4)$$

## Definition 5

If  $\mathbf{x} = (x_1, x_2, \dots, x_n)$  and  $\mathbf{y} = (y_1, y_2, \dots, y_n)$  are two vectors in  $\mathbb{R}^n$ , the  $\ell_2$  and  $\ell_\infty$  distances between  $x$  and  $y$  are defined by

$$\|\mathbf{x} - \mathbf{y}\|_2 = \left( \sum_{i=1}^n |x_i - y_i|^2 \right)^{1/2}$$

and

$$\|\mathbf{x} - \mathbf{y}\|_\infty = \max_{1 \leq i \leq n} |x_i - y_i|$$



# Vector Norms

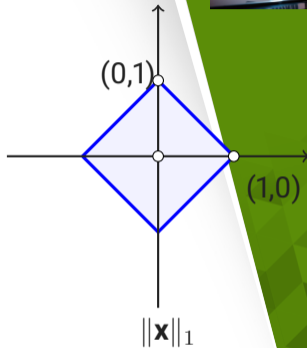
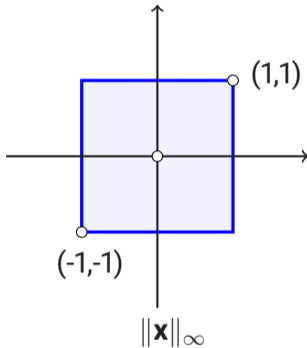
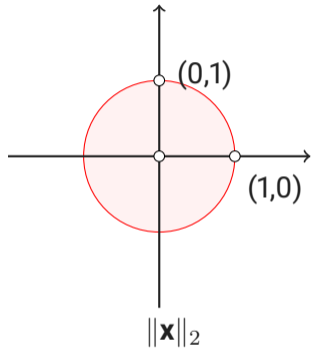


To understand these norms better, see the visualization of the following picture which gives the sketch of the set

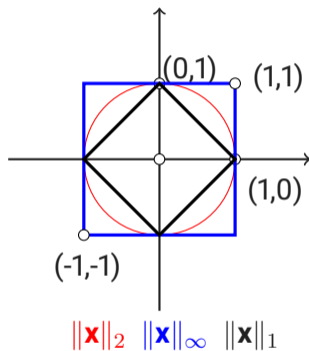
$$\{\mathbf{x} : \mathbf{x} \in \mathbb{R}^2, \|\mathbf{x}\| \leq 1\}$$

This set is called unit cell or the unit ball in two-dimensional vector space.

# Vector Norms



# Vector Norms



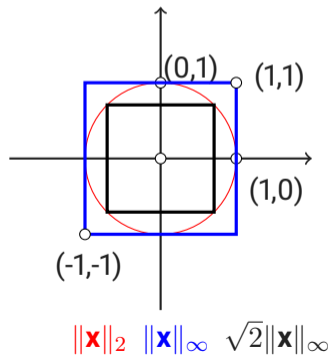
# Vector Norms



## Theorem 6

For each  $x \in \mathbb{R}^n$ ,

$$\|\mathbf{x}\|_{\infty} \leq \|\mathbf{x}\|_2 \leq \sqrt{n}\|\mathbf{x}\|_{\infty}$$



# Matrix Norm



If  $A = (a_{ij})$  is an  $n \times n$  matrix and  $\|\cdot\|$  is a vector norm in  $\mathbb{R}^n$ , then the matrix norm is

$$\|A\| = \max_{\|x\|=1} \|Ax\| \quad (\text{Vector Induced Norm})$$

$$\|A\|_{\infty} = \max_{1 \leq i \leq n} \sum_{j=1}^n |a_{ij}| \quad (\text{Row Sum Norm})$$

$$\|A\|_1 = \max_{1 \leq j \leq n} \sum_{i=1}^n |a_{ij}| \quad (\text{Column Sum Norm})$$

$$\|A\|_F = \sqrt{\sum_{i=1}^n \sum_{j=1}^n |a_{ij}|^2} \quad (\text{Frobenius Norm})$$

$$\|A\|_2 = \sqrt{\rho(A^T A)} = \sigma_{max}(A)$$

# Condition Number



If  $A$  is an invertible matrix, then its condition number  $\kappa(A)$  is defined by

$$\kappa(A) = \|A\| \|A^{-1}\| \quad (5)$$

$$\kappa(A) = \frac{\sigma_{max}(A)}{\sigma_{min}(A)} \quad (l_2 \text{ Norm})$$

$$\kappa(A) = \frac{|\lambda_{max}(A)|}{|\lambda_{min}(A)|} \quad (A * A = AA * A \text{ is normal})$$

$$\kappa(A) = 1 \quad (A * A = AA * = IA \text{ is unitary})$$



# Iterative Methods

# Theorems



## Theorem 7

If  $A$  is an  $n \times n$  matrix such that  $\|A\| < 1$ , then  $I - A$  is invertible, and

$$(I - A)^{-1} = \sum_{k=0}^{\infty} A^k$$

From this theorem, we can observe that

$$\|(I - A)^{-1}\| \leq \sum_{k=0}^{\infty} \|A^k\| \leq \sum_{k=0}^{\infty} \|A\|^k = \frac{1}{1 - \|A\|}$$



# Theorems



## Theorem 8

If  $A$  and  $B$  are  $n \times n$  matrices such that  $\|I - AB\| < 1$ , then  $A$  and  $B$  are invertible. Furthermore,

$$A^{-1} = B \sum_{k=0}^{\infty} (I - AB)^k$$

and

$$B^{-1} = \sum_{k=0}^{\infty} (I - AB)^k A$$

# Iterative Methods

The general algorithm for solving a system  $Ax = b$  is as follows:

1. Choose a nonsingular matrix  $Q$
2. Choose an arbitrary starting vector  $x^{(0)}$
3. Generate vectors  $x^{(1)}, x^{(2)}, \dots$  recursively from the equation

$$Qx^{(k)} = (Q - A)x^{(k-1)} + b, \quad k = 1, 2, 3, \dots \quad (6)$$

Suppose  $x$  is the solution of the system  $Ax = b$  and  $x^{(k)}$  converges to  $x$ , as  $k \rightarrow \infty$ , then

$$Qx = (Q - A)x + b$$

Note, that the system (6) should be easy to solve for  $x^{(k)}$  when the right hand side is known. Also,  $Q$  should be chosen to ensure that  $x^{(k)}$  converges to  $x$ , no matter, what initial vector is used and the convergence should be rapid.

# Iterative Methods

Note that the true solution  $x$  satisfies the equation

$$x = (I - Q^{-1}A)x + Q^{-1}b \quad (7)$$

Therefore,  $x$  is a fixed point of the mapping

$$f(x) = (I - Q^{-1}A)x + Q^{-1}b$$

From (6)

$$\begin{aligned} Qx^{(k)} &= (Q - A)x^{(k-1)} + b, \quad k = 1, 2, 3, \dots \\ \implies x^{(k)} &= (I - Q^{-1}A)x^{(k-1)} + Q^{-1}b, \quad k = 1, 2, 3, \dots \end{aligned}$$

Now,

$$x^{(k)} - x = (I - Q^{-1}A)(x^{(k-1)} - x), \quad k = 1, 2, 3, \dots \quad (8)$$

$$\implies \|x^{(k)} - x\| = \|(I - Q^{-1}A)\| \|x^{(k-1)} - x\|$$

$$\implies \|x^{(k)} - x\| = \|(I - Q^{-1}A)\|^k \|x^{(0)} - x\|$$

# Iterative Methods

If  $\|I - Q^{-1}A\| < 1$ , we can conclude that

$$\lim_{k \rightarrow \infty} \|x^{(k)} - x\| = 0$$

By above theorem, it is guaranteed that if  $\|I - Q^{-1}A\| < 1$ , then both  $Q^{-1}A$  and  $A$  are invertible.

## Theorem 9

If  $\|I - Q^{-1}A\| < 1$  for some matrix norm, then the sequence produced by (6) converges to the solution of  $Ax = b$  for any initial vector  $x^{(0)}$ .

## Theorem 10

If all eigenvalues of  $I - Q^{-1}A$  lies in the open unit disc  $|z| < 1$ , then the sequence produced by (6) converges to the solution of  $Ax = b$  for any initial vector  $x^{(0)}$ .

# Iterative Methods

The above theorem implies that the spectral radius of  $I - Q^{-1}A$  must be less than 1, that is,

$$\rho(I - Q^{-1}A) < 1$$

Let  $r^{(k)}$  denote the residual vector obtained from  $x^{(k)}$  after  $k$  iterations, then we get

$$r^{(k)} = b - Ax^{(k)}.$$

By above theorem, if  $\|I - Q^{-1}A\| < 1$ , then  $\|r^{(k)}\| \rightarrow 0$ .

# Theorem

The distance between two matrices  $A, B \in M_{n \times n}(\mathbb{R})$  is  $\|A - B\|$ .

## Theorem 11

If  $\|\cdot\|$  is a vector norm in  $\mathbb{R}^n$ , then the matrix norm is

$$\|A\| = \max_{\|x\|=1} \|Ax\|$$

# Theorems and Proof



## Theorem 12

If  $\|\cdot\|$  is a vector norm in  $\mathbb{R}^n$ , then the matrix norm is

$$\|A\| = \max_{\|x\|=1} \|Ax\|$$

**Proof:** It is enough to prove that  $\|\cdot\|$  is a matrix norm. If  $A \neq 0$ , at least one of its column is not a zero vector. Let  $j$ th column  $(A)_j \neq 0$ . Then

$$\|A\| = \max_{\|x\|=1} \|Ax\| \geq \|Ae_j\| = \|(A)_j\| > 0$$

$$\|\alpha A\| = \max_{\|x\|=1} \|\alpha Ax\| = \max_{\|x\|=1} |\alpha| \|Ax\| = |\alpha| \max_{\|x\|=1} \|Ax\| = |\alpha| \|A\|$$

# Theorems and Proof



Now,

$$\|A + B\| = \max_{\|x\|=1} \|(A + B)x\| = \max_{\|x\|=1} \|Ax + Bx\| \leq \max_{\|x\|=1} (\|Ax\| + \|Bx\|)$$

$$\max_{\|x\|=1} \|Ax\| + \max_{\|x\|=1} \|Bx\| = \|A\| + \|B\|$$

$$\|AB\| = \max_{\|x\|=1} \|ABx\| \leq \max_{\|x\|=1} \|A\| \|Bx\| \leq \max_{\|x\|=1} \|A\| \|B\| \|x\| = \|A\| \|B\|$$

Hence the proof.



# Theorem



## Theorem 13

If  $A = (a_{ij})$  is an  $n \times n$  matrix, then

$$\|A\|_{\infty} = \max_{1 \leq i \leq n} \sum_{j=1}^n |a_{ij}|$$

$$\|A\|_1 = \max_{1 \leq j \leq n} \sum_{i=1}^n |a_{ij}|$$

The  $\|A\|_1$  norm is also called as the column sum norm as it computes the maximum absolute column sum of the matrix. The  $\|A\|_{\infty}$  norm is also called as the row sum norm as it computes the maximum absolute row sum of the matrix.

# Theorems and Proof

## Theorem 14

If  $A = (a_{ij})$  is an  $n \times n$  matrix, then

$$\|A\|_{\infty} = \max_{1 \leq i \leq n} \sum_{j=1}^n |a_{ij}|$$

$$\|A\|_1 = \max_{1 \leq j \leq n} \sum_{i=1}^n |a_{ij}|$$

**Proof:** Observe that when  $Ax = b$

$$b_i = \sum_{j=1}^n a_{ij}x_j \implies \|b\|_{\infty} = \max_{1 \leq i \leq n} |b_i| = \max_{1 \leq i \leq n} \left| \sum_{j=1}^n a_{ij}x_j \right|$$

# Theorems and Proof



$$\|A\|_{\infty} = \max_{\|x\|_{\infty}=1} \|Ax\| = \max_{\|x\|_{\infty}=1} \max_{1 \leq i \leq n} \sum_{j=1}^n |a_{ij}x_j| \leq \max_{\|x\|_{\infty}=1} \max_{1 \leq i \leq n} \sum_{j=1}^n |a_{ij}||x_j|$$

$$\Rightarrow \|A\|_{\infty} \leq \left( \max_{1 \leq i \leq n} \sum_{j=1}^n |a_{ij}| \right) \max_{\|x\|_{\infty}=1} |x_j| = \left( \max_{1 \leq i \leq n} \sum_{j=1}^n |a_{ij}| \right) \max_{\|x\|_{\infty}=1} \|x\|$$

$$\Rightarrow \|A\|_{\infty} \leq \max_{1 \leq i \leq n} \sum_{j=1}^n |a_{ij}|$$

It is enough to prove that the equality is achieved for some  $x$ . Choose your  $x$  such that  $x_j = \text{sign}(A_{ij})$ .

# Theorems and Proof

Then the proof follows for row sum follows

$$\|A\|_{\infty} = \max_{1 \leq i \leq n} \sum_{j=1}^n |a_{ij}|$$

For column sum, consider the columns of  $A$  as  $A_1, A_2, \dots, A_n$

$$\|Ax\|_1 = \left\| \sum_{i=1}^n A_i x \right\|_1 \leq \sum_{i=1}^n \|A_i x\|_1 = \sum_{i=1}^n |x| \|A_i\|_1 \leq \max_{1 \leq j \leq n} \|x\|_1 \sum_{i=1}^n |a_{ij}|$$

$$\implies \|A\|_1 = \max_{\|x\|_1=1} \|Ax\|_1 \leq \max_{1 \leq j \leq n} \sum_{i=1}^n |a_{ij}|$$

The equality and the proof follows by taking the standard basis vector for  $x$ .

# Theorem

The Frobenius norm is given by  $\|A\|_F$

$$\|A\|_F = \sqrt{\sum_{i=1}^n \sum_{j=1}^n |a_{ij}|^2}$$



# Theorems and Proof



## Theorem 15

$$\sum_{i=1}^n \sum_{j=1}^n a_{ij}^2 = \mathbf{trace}(A^T A)$$

$$\|A\|_F^2 = \mathbf{trace}(A^T A)$$

**Proof:** If  $A = (a_{ij})$  and  $B = (b_{ij})$ , then  $C = AB = (c_{ij})$  is given by, then

$$c_{ij} = \sum_{k=1}^n a_{ik} b_{kj} \implies c_{ii} = \sum_{k=1}^n a_{ik} b_{ki}$$

# Theorems and Proof

When  $C = A^T A$ , we have

$$c_{ii} = \sum_{k=1}^n a_{ki} a_{ki} = \sum_{k=1}^n a_{ki}^2$$

$$\text{trace}(A^T A) = \sum_{i=1}^n c_{ii} = \sum_{i=1}^n \sum_{k=1}^n a_{ki}^2 = \sum_{j=1}^n \sum_{i=1}^n a_{ij}^2 = \sum_{i=1}^n \sum_{j=1}^n a_{ij}^2$$

Since

$$\sum_{i=1}^n \sum_{j=1}^n a_{ij}^2 = \sum_{i=1}^n \sum_{j=1}^n |a_{ij}|^2$$

the proof follows.

# Theorems and Proof



## Theorem 16

Frobenius norm is a norm.

**Proof:** If  $A \neq 0$ , at least one element is not a zero element. Let it be  $a_{pk}$ .  
Hence

$$\|A\|_F = \sqrt{\sum_{i=1}^n \sum_{j=1}^n |a_{ij}|^2} \geq |a_{pk}| > 0$$

$$\|\alpha A\|_F = \sqrt{\sum_{i=1}^n \sum_{j=1}^n |\alpha a_{ij}|^2} = \sqrt{\sum_{i=1}^n \sum_{j=1}^n |\alpha|^2 |a_{ij}|^2} = |\alpha| \sqrt{\sum_{i=1}^n \sum_{j=1}^n |a_{ij}|^2} = |\alpha| \|A\|_F$$



# Theorems and Proof



$$\|A + B\|_F^2 = \sum_{i=1}^n \sum_{j=1}^n |a_{ij} + b_{ij}|^2$$

We know

$$|a + b|^2 \leq |a|^2 + |b|^2 + 2|a||b|$$

Hence

$$\|A + B\|_F^2 = \sum_{i=1}^n \sum_{j=1}^n |a_{ij}|^2 + \sum_{i=1}^n \sum_{j=1}^n |b_{ij}|^2 + 2 \sum_{i=1}^n \sum_{j=1}^n |a_{ij}||b_{ij}|$$

$$\|A + B\|_F^2 = \|A\|_F^2 + \|B\|_F^2 + 2 \sum_{i=1}^n \sum_{j=1}^n |a_{ij}||b_{ij}|$$

# Theorems and Proof



By Cauchy-Schwarz inequality

$$\sum_{i=1}^n \sum_{j=1}^n |a_{ij}| |b_{ij}| \leq \left( \sum_{i=1}^n \sum_{j=1}^n |a_{ij}|^2 \right)^{1/2} \left( \sum_{i=1}^n \sum_{j=1}^n |b_{ij}|^2 \right)^{1/2} = \|A\|_F \|B\|_F$$

$$\|A + B\|_F^2 = \|A\|_F^2 + \|B\|_F^2 + 2\|A\|_F \|B\|_F \implies \|A + B\|_F \leq \|A\|_F + \|B\|_F$$

$$\|AB\|_F^2 = \sum_{i=1}^n \sum_{j=1}^n \left| \sum_{k=1}^n a_{ik} b_{kj} \right|^2 \leq \sum_{i=1}^n \sum_{j=1}^n \left( \sum_{k=1}^n |a_{ik}|^2 \sum_{k=1}^n |b_{kj}|^2 \right)$$

$$\|AB\|_F^2 \leq \sum_{i=1}^n \sum_{j=1}^n \left( \sum_{k=1}^n |a_{ik}|^2 \sum_{l=1}^n |b_{lj}|^2 \right) = \sum_{i=1}^n \sum_{k=1}^n |a_{ik}|^2 \sum_{l=1}^n \sum_{j=1}^n |b_{lj}|^2 = \|A\|_F^2 \|B\|_F^2$$

# Theorems and Proof

## Theorem 17

If  $U$  and  $V$  are orthogonal, then

$$\|UA\|_F = \|AV\|_F = \|A\|_F$$

**Proof:**

$$\|UA\|_F^2 = \mathbf{trace}((UA)^T(UA)) = \mathbf{trace}(A^T U^T U A) = \mathbf{trace}(A^T A) = \|A\|_F^2$$

Since  $\mathbf{trace}(A^T A) = \mathbf{trace}(AA^T)$ , we get

$$\begin{aligned} \|AV\|_F^2 &= \mathbf{trace}((AV)^T(AV)) = \mathbf{trace}((AV)(AV)^T) \\ &= \mathbf{trace}(AVV^T A^T) = \mathbf{trace}(AA^T) = \mathbf{trace}(A^T A) = \|A\|_F^2 \end{aligned}$$

# Example

## Example 18

Compute the  $\|\cdot\|_1$ ,  $\|\cdot\|_F$  and  $\|\cdot\|_\infty$  norms of the following matrix.

$$\begin{pmatrix} 6 & -2 & 2 & 4 \\ 12 & -8 & 6 & 10 \\ 3 & -13 & 9 & 3 \\ -6 & 4 & 1 & -18 \end{pmatrix}$$

$$\|A\|_1 = \max_{1 \leq j \leq n} \sum_{i=1}^n |a_{ij}| = \max\{27, 27, 18, 35\} = 35$$

$$\|A\|_\infty = \max_{1 \leq i \leq n} \sum_{j=1}^n |a_{ij}| = \max\{14, 36, 28, 29\} = 36$$

$$\|A\|_F = 32.3883$$

# Eigenvalues

If  $I$  denotes the identity matrix, then  $\|I\| = 1$ . Also,

$$\|AB\| \leq \|A\|\|B\|$$

## Definition 19 (Characteristic Polynomial:)

If  $A$  is a square matrix, the characteristic polynomial of  $A$  is defined by

$$p(\lambda) = \det(A - \lambda I)$$

## Definition 20 (Eigenvalue and Eigenvector)

If  $p$  is a characteristic polynomial of  $A$ , the roots of  $p$  are called eigenvalues or characteristic values of  $A$ . If  $\lambda$  is an eigenvalue of  $A$  and  $x \neq 0$  such that  $Ax = \lambda x$ , then  $x$  is an eigenvector of  $A$ .

# Spectral Radius

Usually, the matrix is considered over the field  $C$  though the matrix entries can be real.

## Definition 21 (Singular Values)

The singular values  $\sigma$  of an  $m \times n$  matrix  $A$  are the positive square roots of the nonzero eigenvalues of the  $n \times n$  symmetric matrix  $A^T A$ .

$$\sigma_i(A) = \sqrt{\lambda_i(A^T A)}$$

## Definition 22 (Spectral Radius)

The spectral radius  $\rho(A)$  of a matrix  $A$  is defined by

$$\rho(A) = \max_{1 \leq i \leq n} |\lambda_i|$$

This is also denoted as  $\lambda_{max}(A)$ .

# Spectral Radius

The matrix norm induced by the Euclidean norm ( $\ell_2$  norm) is the spectral norm, and is given by

$$\|A\|_2 = \sqrt{\rho(A^T A)} = \sigma_{max}(A)$$

This is the minimum norm and provides the tightest measure of size.

## Theorem 23

If  $A = (a_{ij})$  is an  $n \times n$  matrix, then

$$\|A\|_2 \leq \|A\|_F \leq \sqrt{n} \|A\|_2,$$

$$\|A\|_F = \sqrt{\text{trace}(A^T A)}$$

$$\rho(A) \leq \|A\|$$

for any natural norm  $\|\cdot\|$ .

# Thanks

**Doubts and Suggestions**

[panch.m@iittp.ac.in](mailto:panch.m@iittp.ac.in)





# MA633L-Numerical Analysis

Lecture 26 : Numerical Linear Algebra - Norms

**Panchatcharam Mariappan<sup>1</sup>**

<sup>1</sup>Associate Professor  
Department of Mathematics and Statistics  
IIT Tirupati, Tirupati

**March 10, 2025**

